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PGX.D aDFS: An Almost Depth-First-Search Distributed Graph-Querying System

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Complexities in Graph-Query Execution

- Limited locality (especially in a distributed system)
- Intermediate (and final) result explosion

Need a distributed solution that is flexible and can handle the scale

Twitter graph

```
SELECT COUNT(*) MATCH (a)->()
+----------------------+
|    COUNT(*)          |
+----------------------+
| 1,468,365,182        |
+----------------------+
```

1 hop

```
SELECT COUNT(*) MATCH (a)->()->()
+----------------------+
|    COUNT(*)          |
+----------------------+
| 9,324,563,362,739    |
+----------------------+
```

2 hops

spoiler!
PGX.D aDFS
8 machines
~20 minutes
~8B matches/s

```
-- Info of authors who like each other and have < 10 years of age difference
SELECT a1.name, a2.name, a1.country = a2.country,
   ABS(a1.salary - a2.salary) AS salary_diff
MATCH (a1:author) -[:likes]-> (a2:author) -[:likes]-> (a1)
WHERE ABS(a1.age - a2.age) < 10
ORDER BY salary_diff DESC
```

Any user expression in projections and filters

Requires **homomorphic matching** and returns **all result permutations**

Agenda

1. Introduction / Motivation
2. aDFS Design
3. Evaluation
4. Conclusions

aDFS Design Principles

1. Asynchronous operation

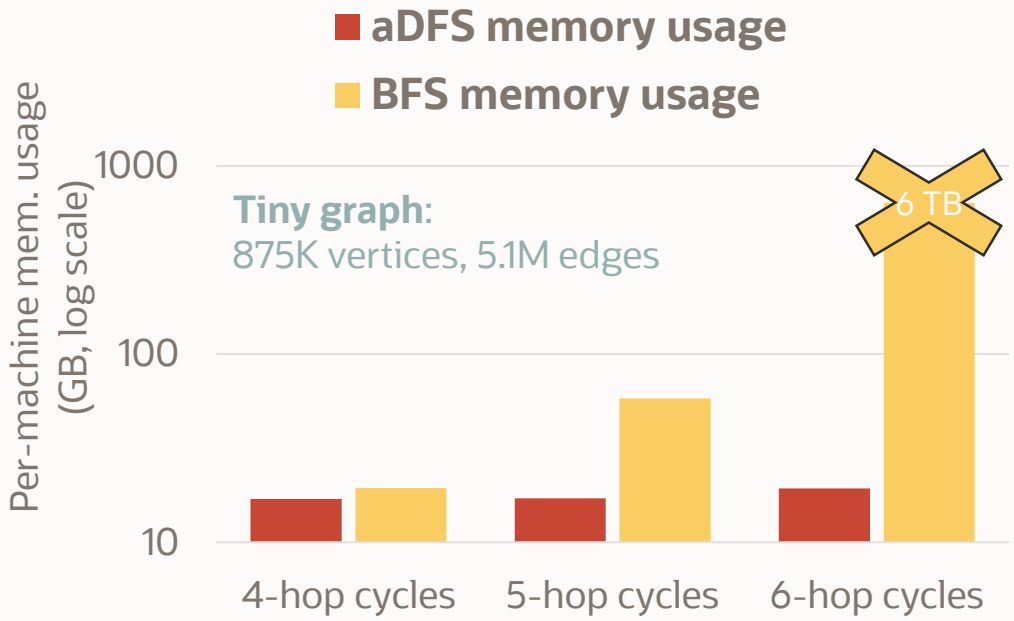
- Workers operate independently
  - on traversals where there is work
- Workers buffer and forget remote traversals

→Workers do not block due to remote communication

2. (Almost) Depth-first traversal

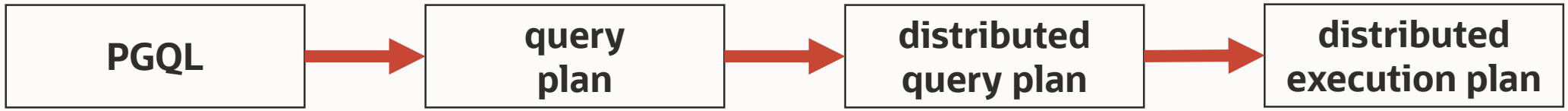
- Eager completion of matches
- Fine-grained flow control

→ Control memory consumption

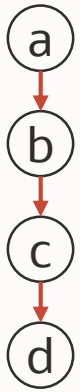


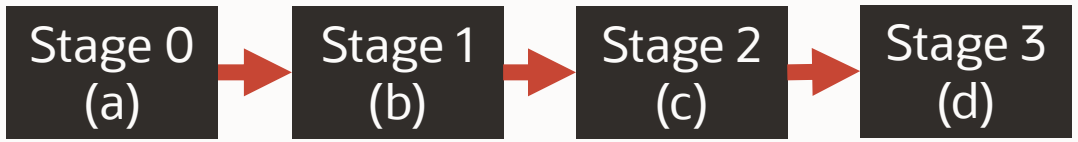
In-memory distributed execution with controllable memory usage

From a PGQL Query to an aDFS Execution Plan




```
SELECT ...
MATCH
  (a)->(b),
  (b)->(c),
  (c)->(d)
WHERE
  ...
```





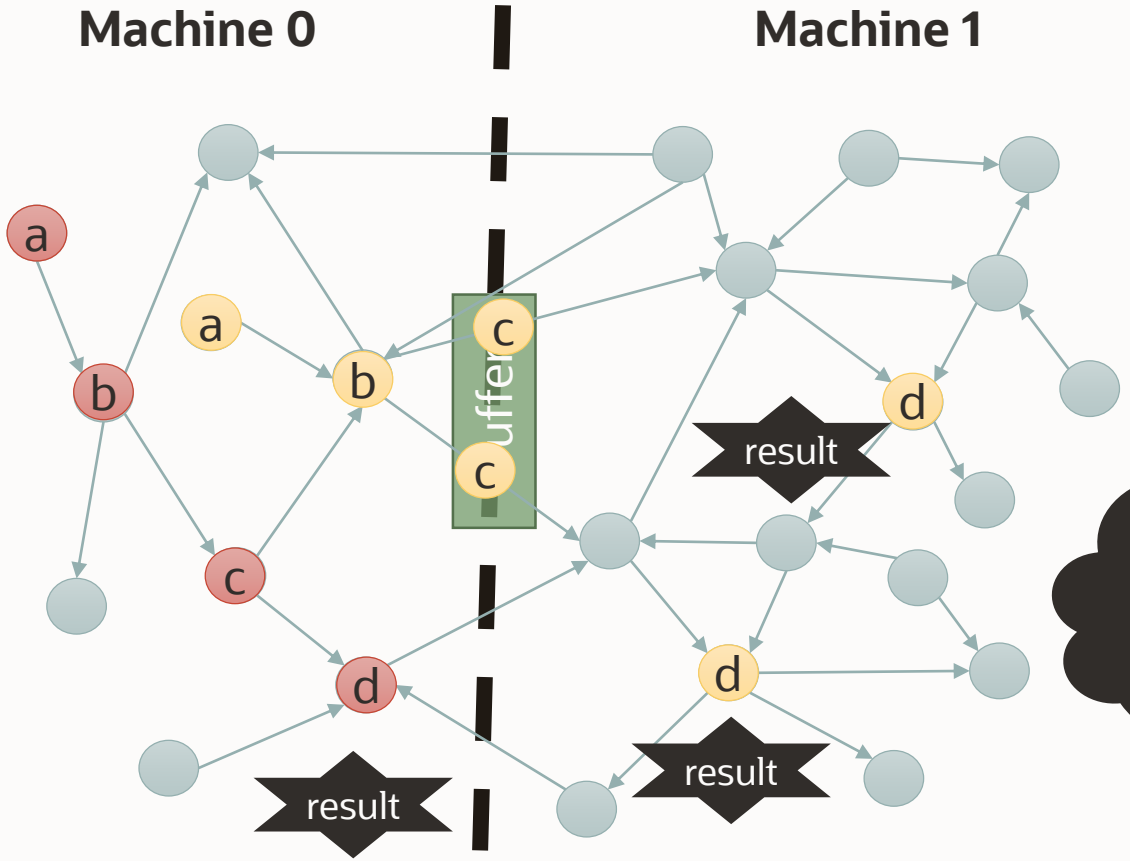
A list of stages that "know" how to

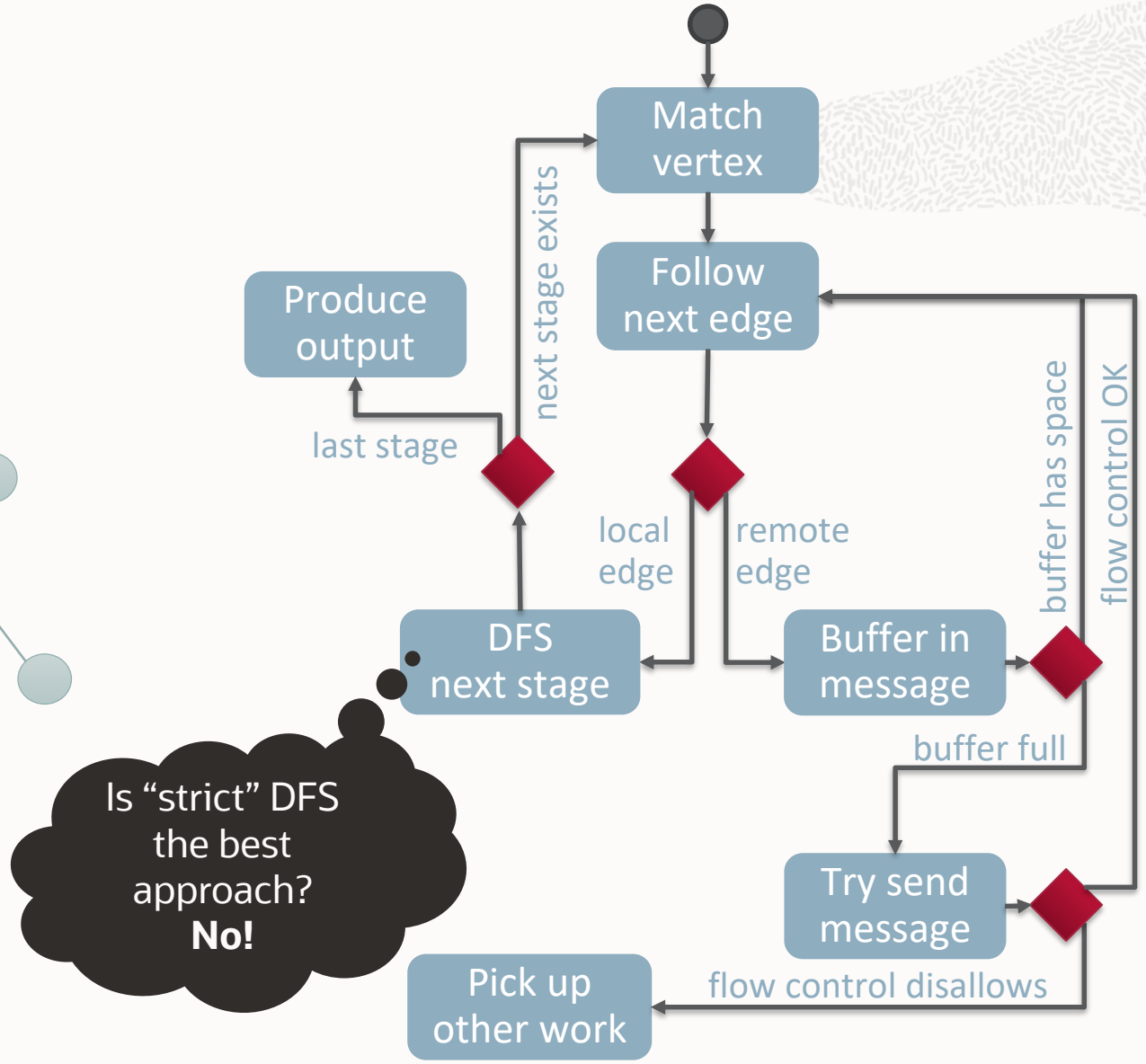
1. match a vertex
2. move to next stage

Asynchronous DFS/BFS Traversals

stage 0 stage 1 stage 2 stage 3

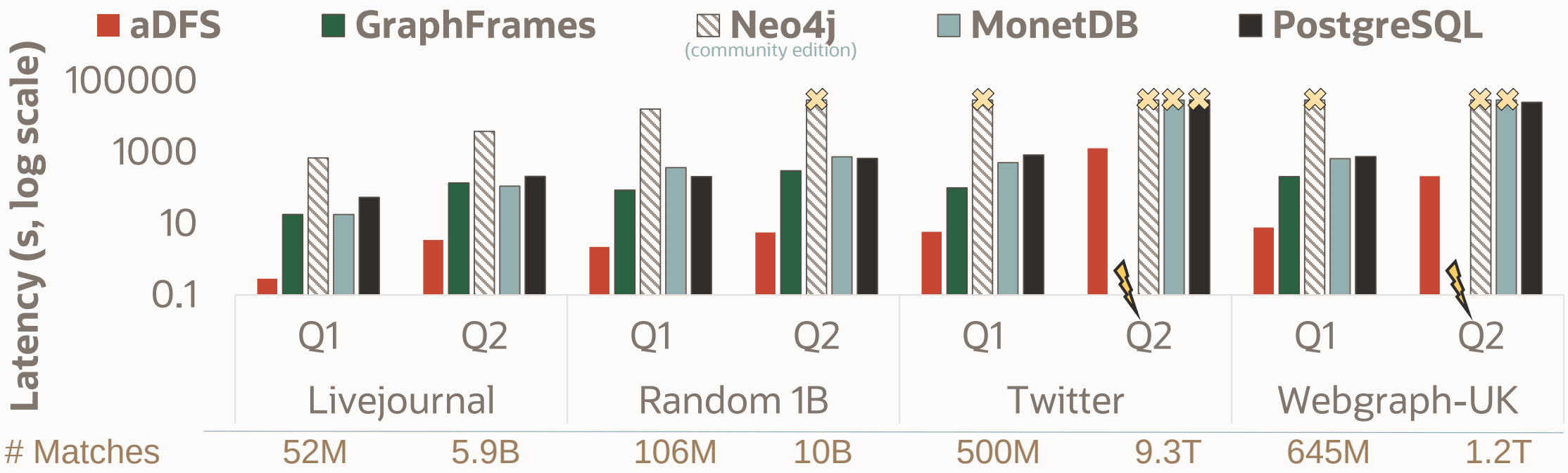
MATCH (a) → (b) → (c) → (d)





Experimental Evaluation

Schemaless Graphs and Queries



- **Q1: cycle (a)->(b)->(a) Q2: 2-hops (a)->(b)->(c)**
- aDFS and GraphFrames with 8 machines / others single machine
- aDFS configured with 1GB memory per machine / others have access to whole machine memory (768 GB)

Did not complete in 8 hours Hang due to out of memory

Only aDFS can handle the scale

Conclusions

- **aDFS is a fast and scalable distributed graph querying engine**
  - Provides flexible PGQL querying
  - Combines BFS / DFS
  - Limits max memory usage

- **Also in the paper:** Experiments with the LDBC graph and queries

- **Since the experiments for ATC paper, PGX.D**
  - supports graphs with schema → lower memory and better performance
  - has significantly faster PGQL query execution
  - supports bigger subset of PGQL

Thank you!