# Evaluating Cypher queries as algebraic expressions within RedisGraph

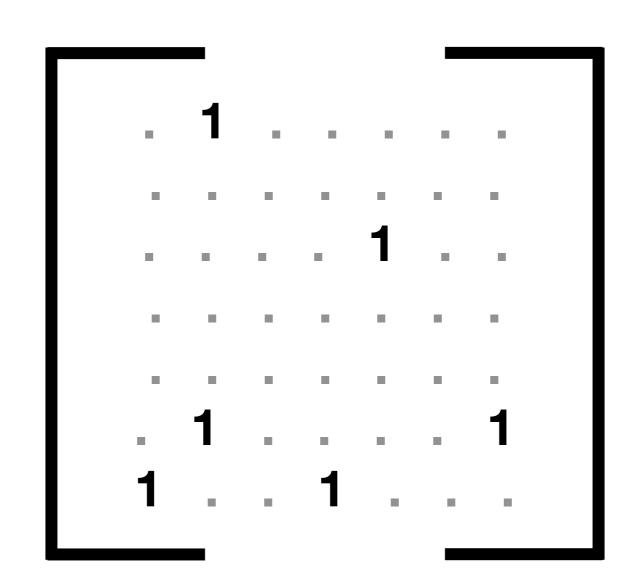
Roi Lipman RedisLabs

# RedisGraph property graph representation

# 4 types of matrices

- THE adjacency matrix
- Label matrices
- Relation matrices
- Relationship mapping matrices

All matrices share the same dimensions #rows = #columns = #nodes



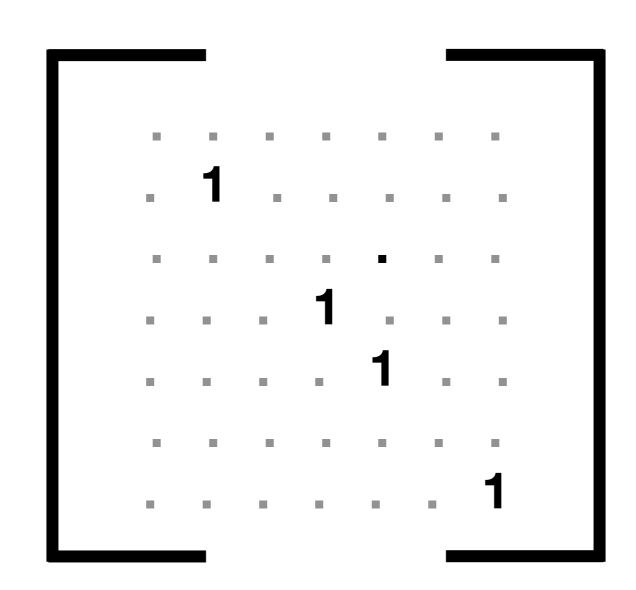
Boolean

- Boolean
- Relation type agnostic

- Boolean
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- Label agnostic

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- Label agnostic
- Directional, M[i, j] = 1

Node i connected to Node j
\*might be multiple times with different
relationship types



Boolean

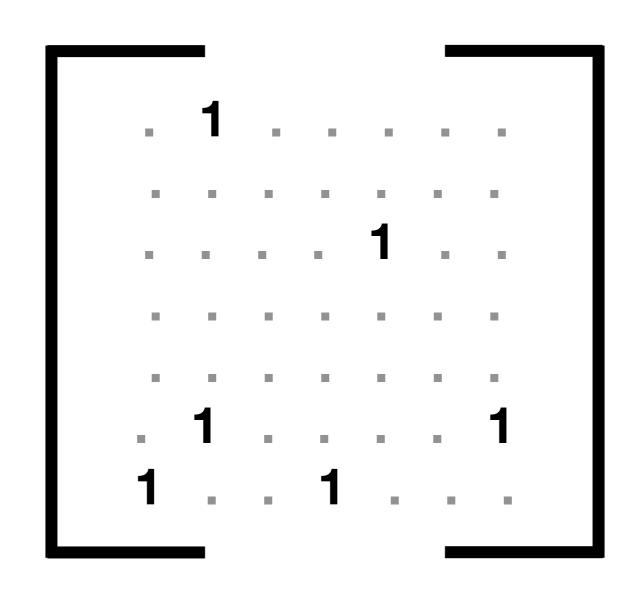
- Boolean
- One for each node type

- Boolean
- One for each node type
- Diagonal

$$L[i, i] = 1$$

Node with ID i is labeled as L

### Relation matrix

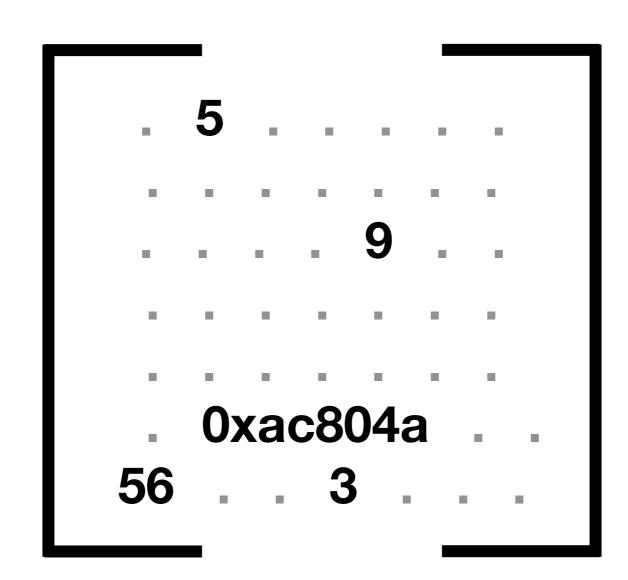


#### Relation matrix

- Boolean
- Label agnostic
- Directional, R[i, j] = 1

Node i connected to Node j \*might be multiple times with relationship type R

# Relation mapping matrix



# Relation mapping matrix

64bit entries

RM[i, j] - either edge ID if node i is connected to node j with a single edge of type R

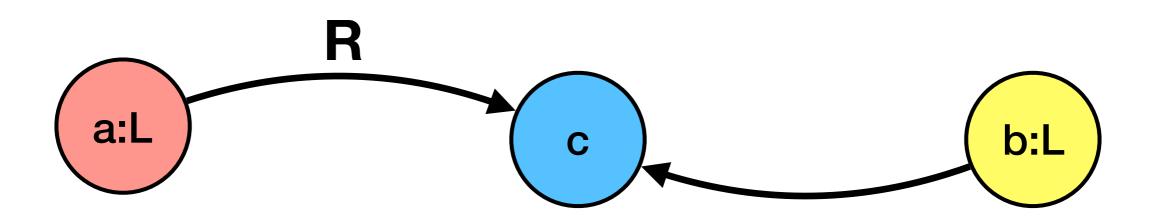
Pointer to edge IDs array if node i is connected to node j with multiple edges of type R

# Query Execution Traversals

#### Chain:

A path where each node on the path appears only once and the sum of its In/Out degrees < 3

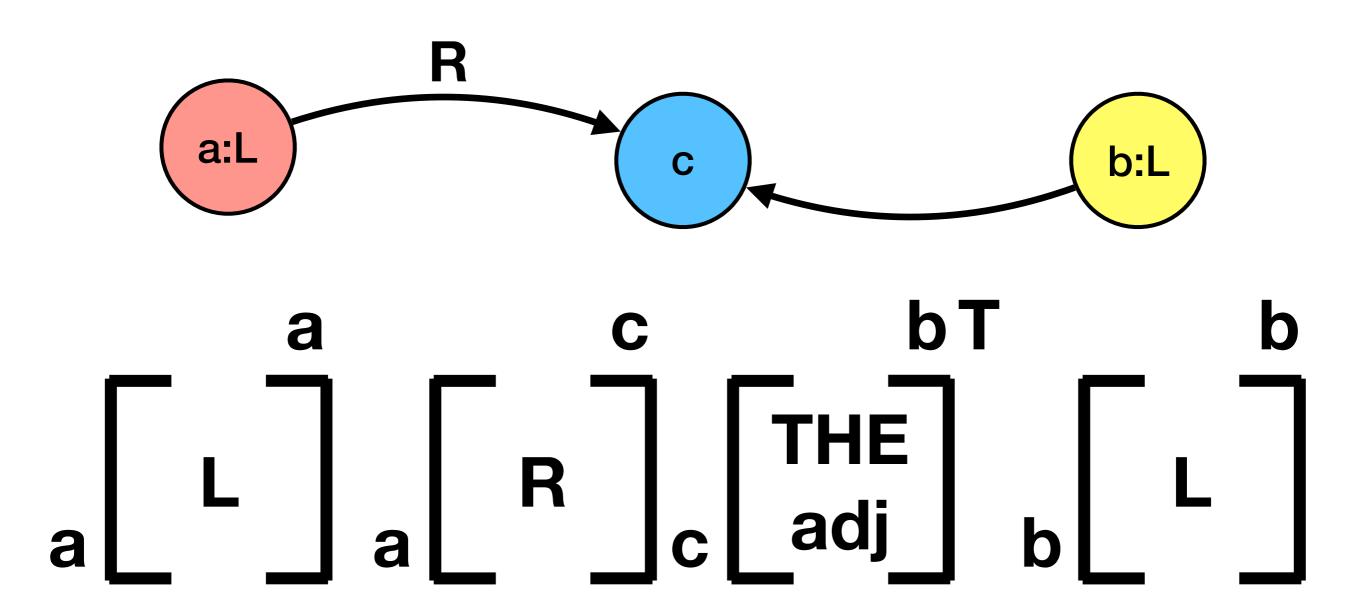
# Query Execution Traversals



MATCH (a:L)-[:R]->(c)<-[]-(b:L)

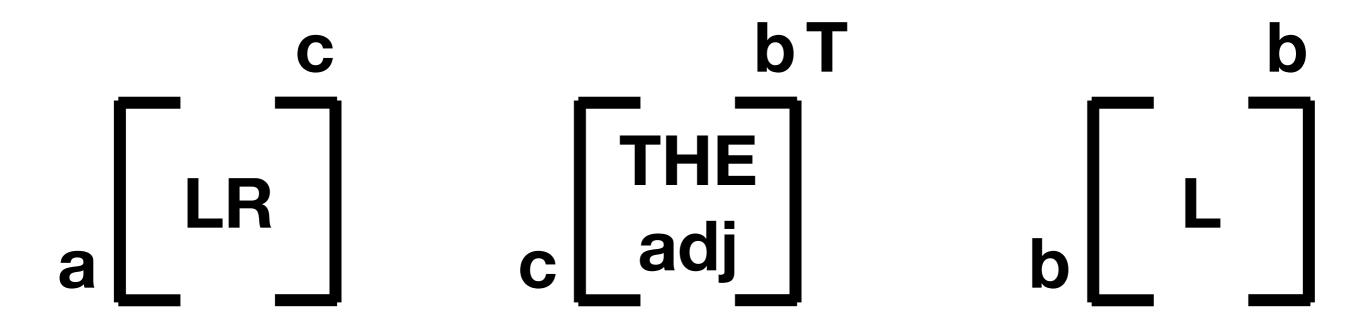
RETURN b

# Algebraic expression



$$\begin{bmatrix} a & c & bT & b\\ a & \begin{bmatrix} C & THE \\ adj \end{bmatrix} & \begin{bmatrix} b \\ L \end{bmatrix} \end{bmatrix}$$

Matrix multiplication is associative Motivation keep sparsity



$$\begin{bmatrix} \mathbf{b} & \mathbf{c} & \mathbf{c} \\ \mathbf{x} \end{bmatrix} * \mathbf{b} \begin{bmatrix} \mathbf{Y} \end{bmatrix} = \mathbf{a} \begin{bmatrix} \mathbf{z} \end{bmatrix}$$

Row domain is "Sticky"

$$\begin{bmatrix} b & c & c \\ x & x & z \end{bmatrix} * \begin{bmatrix} c & c & c \\ c & z & z \end{bmatrix}$$

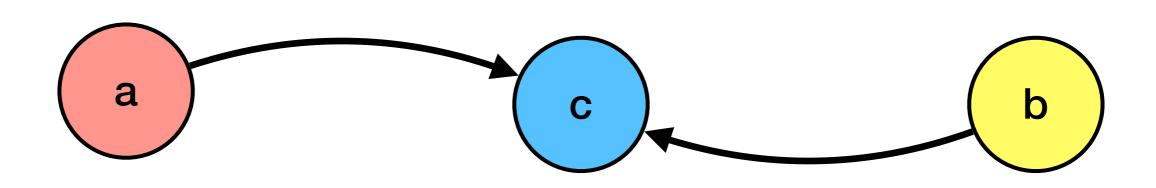
Column domain changes with every "hop"

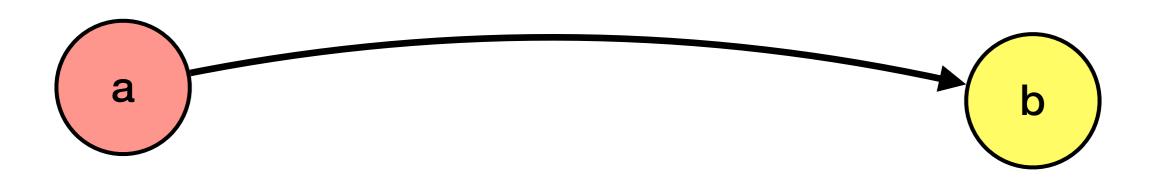
## "2D"

## MxM Lose of information

$$\begin{bmatrix} b & c & c \\ x & x & z \end{bmatrix} * \begin{bmatrix} c & c & c \\ c & z & z \end{bmatrix}$$

# a is connected to b But how?





# MATCH (a:L)-[:R]->(c)<-[]-(b:L) RETURN c

$$\begin{bmatrix} a & c & bT & b\\ L & a & R & THE \\ a & a & c & adj & b \end{bmatrix}$$

#### Intermediate entities

Referred node/edge

RETURN n,e

#### Intermediate entities

Variable length edges

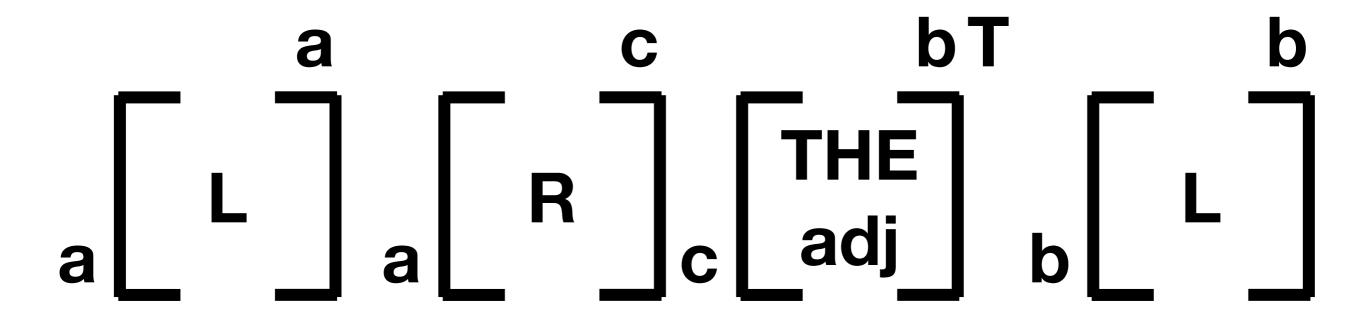
MATCH (a)-[e:\*2..4]->(b)

### Intermediate entities

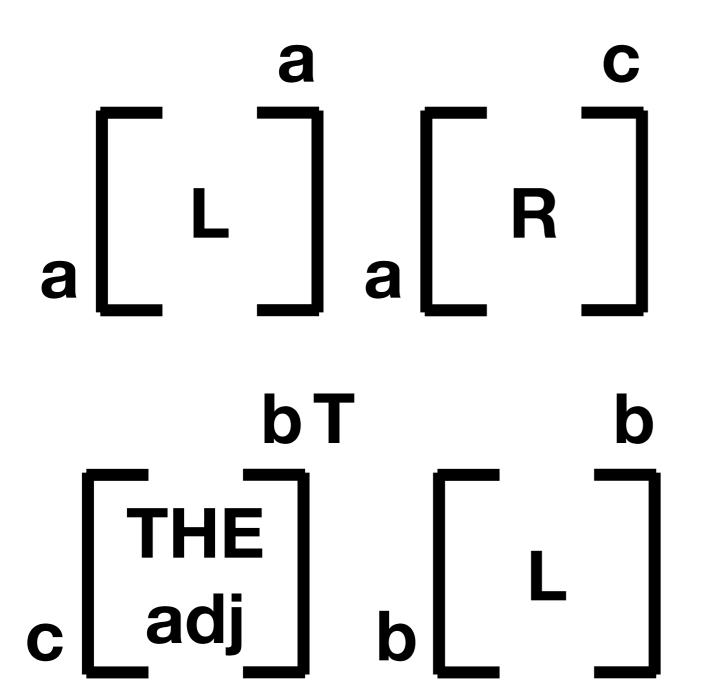
Filtered entities

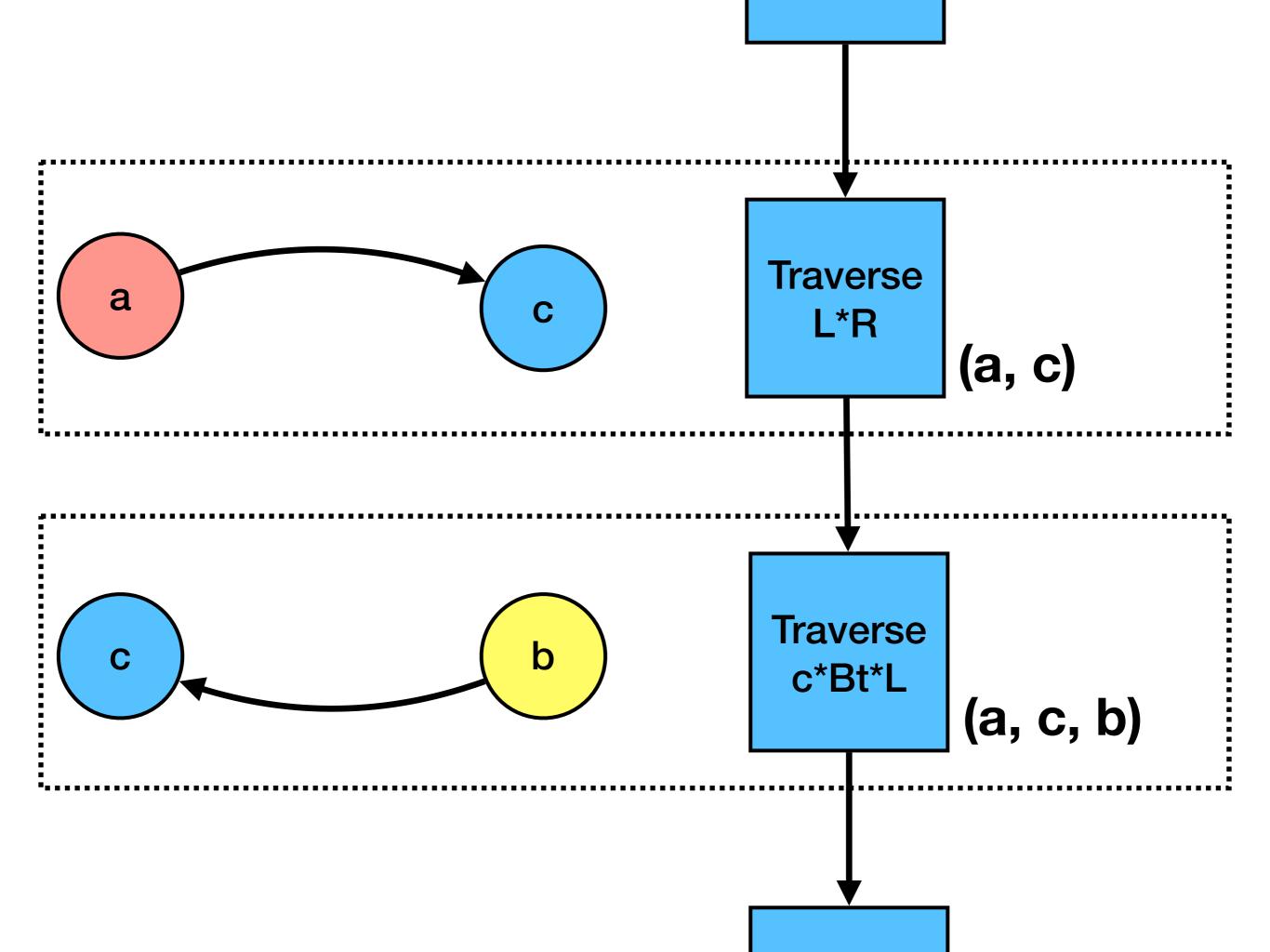
WHERE n.v = 34

#### Break on intermediate

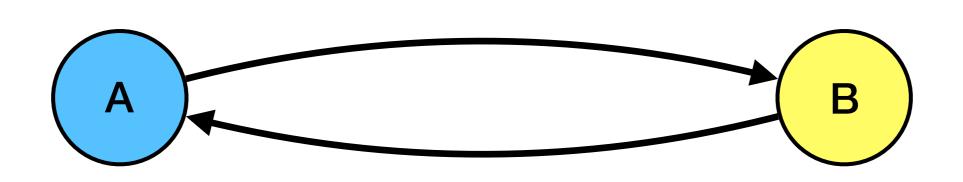


#### Break on intermediate



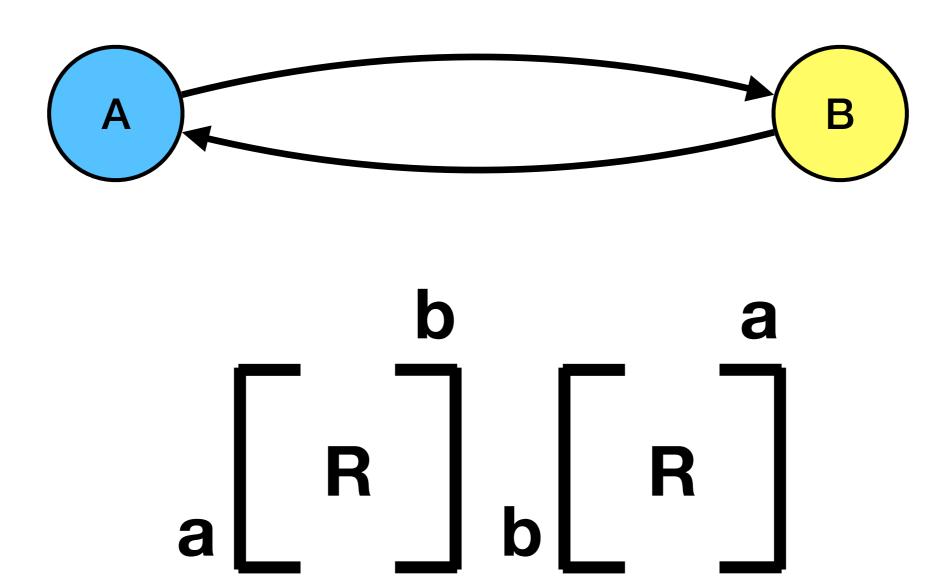


# High degree nodes & Cycles

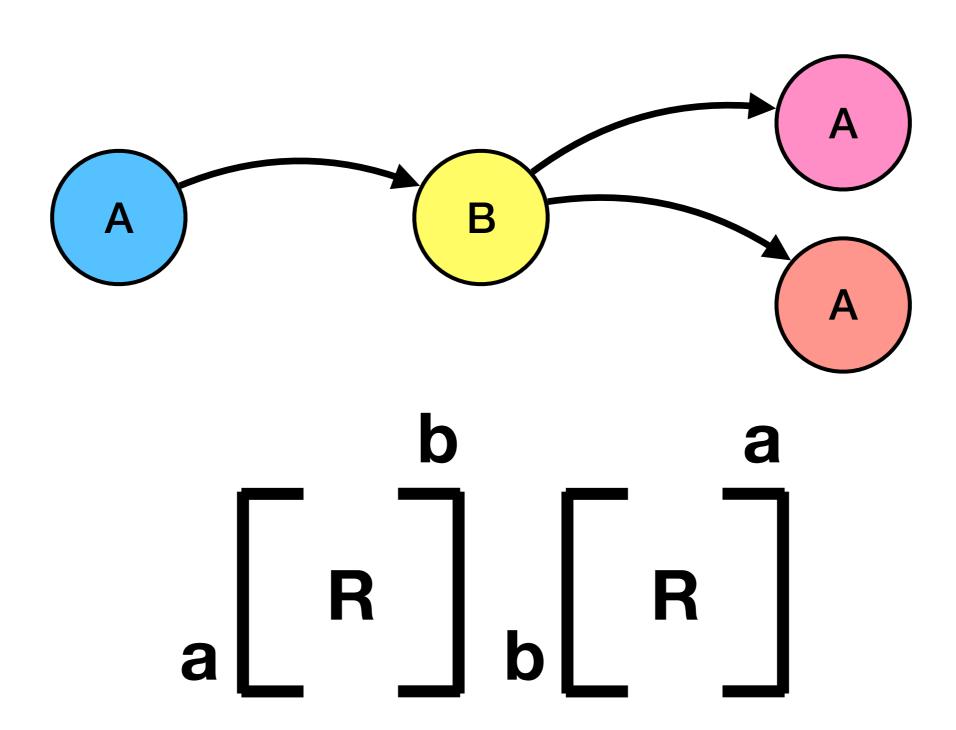


MATCH (a)-[]->(b)-[]->(a)
RETURN a

# Cycle



# Cycle

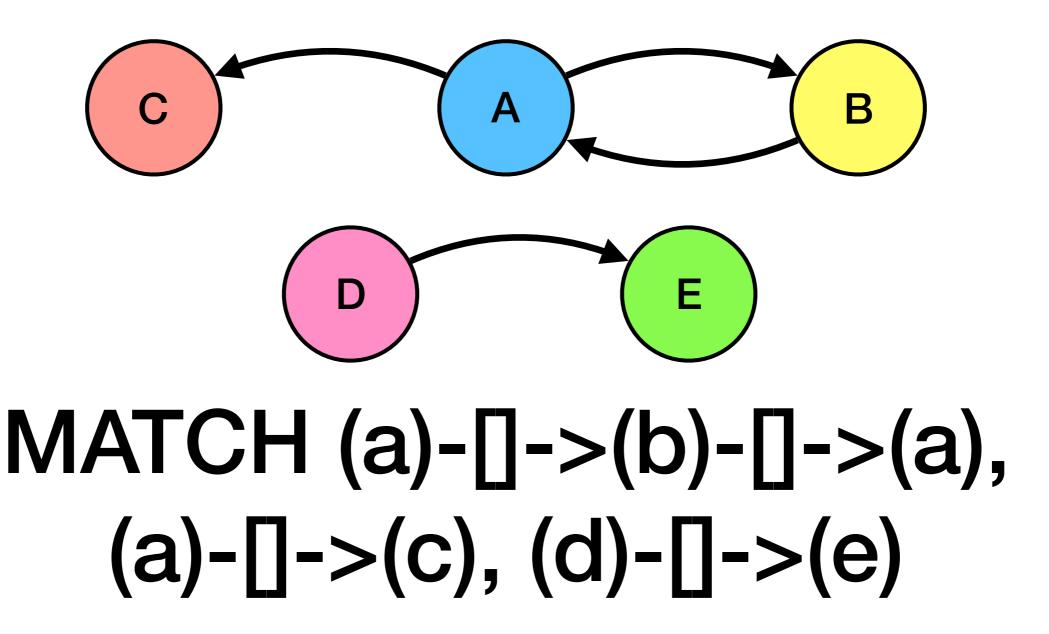


# Improved Algebraic expressions construction

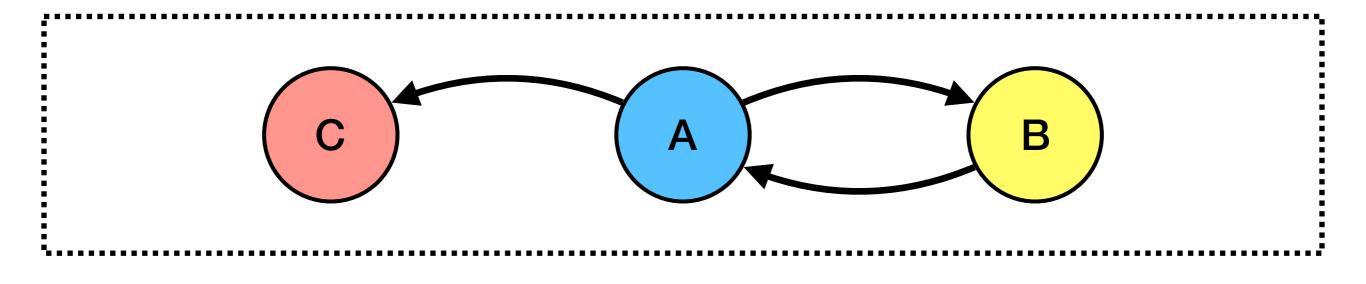
```
Exps = []
QG = QueryGraph(AST)
CCS = ConnectedComponents(QG)
For CC in CCS
    while(CC not empty)
        P = LongestPath(CC)
        E = AlgebricExpressionFromPath(P)
        Es = AlgebricExpressionBreakInter(E)
        Exps += Es
        CC -= P
```

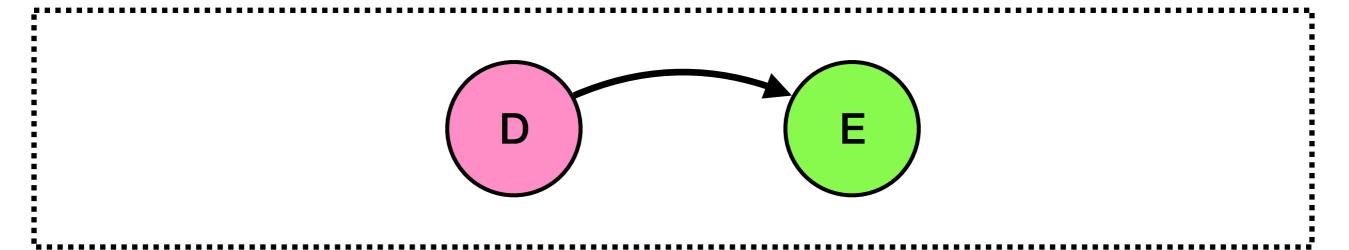
Pseudo code

#### Query Graph

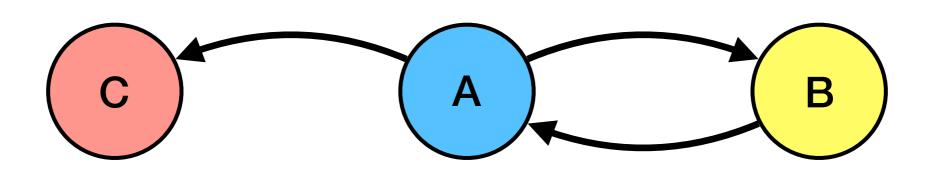


### Connected Components





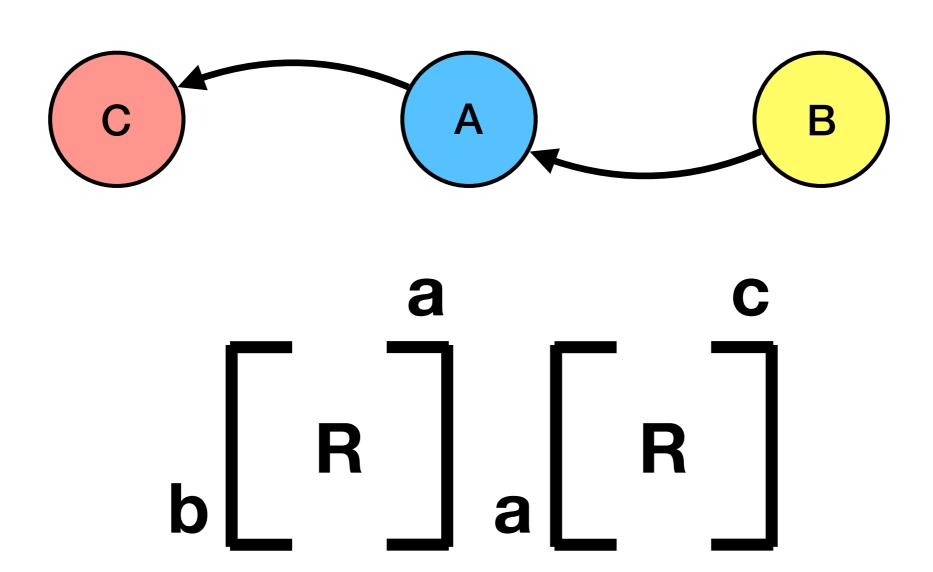
#### Longest path



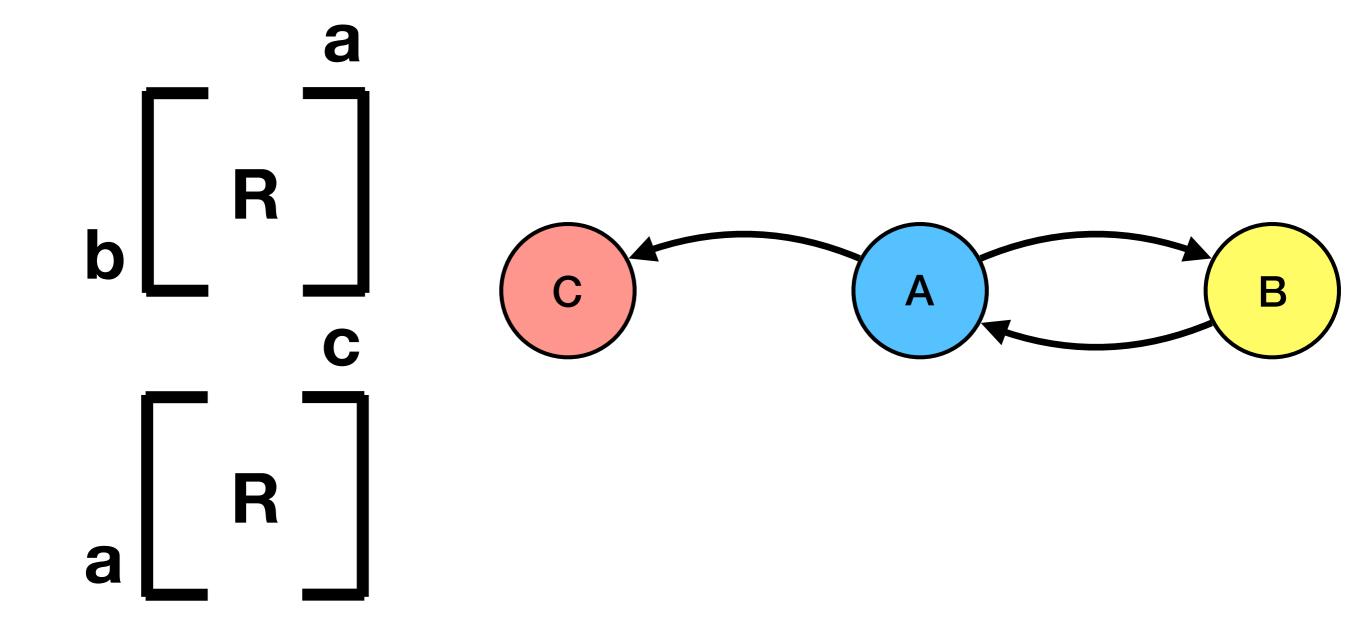
$$P0 = \{C, A, B\}$$

$$P1 = \{B, A, C\}$$

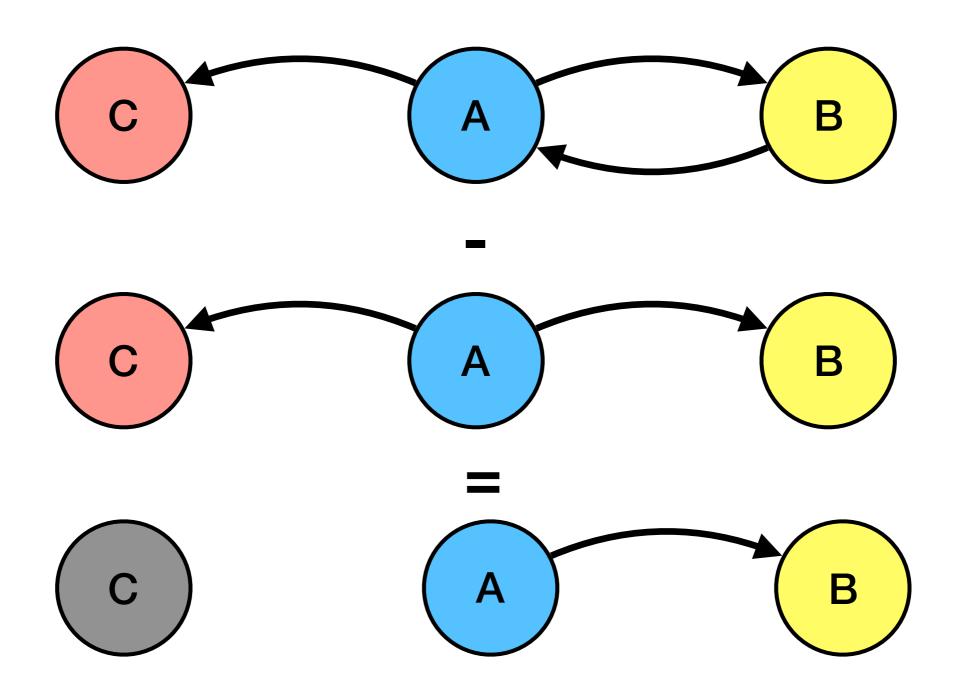
### Algebraic expression from path (chain)



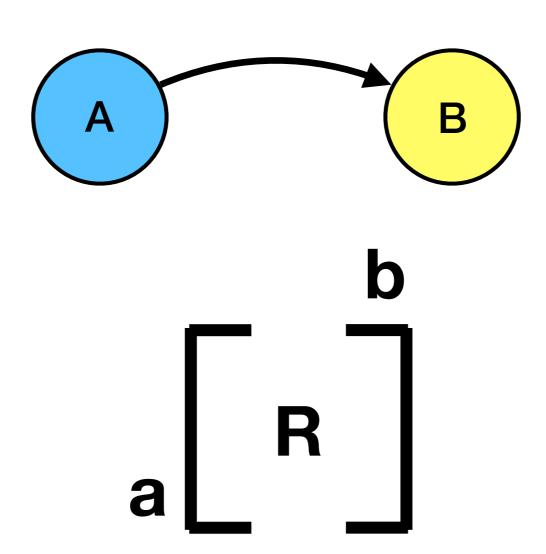
#### Break on intermediates



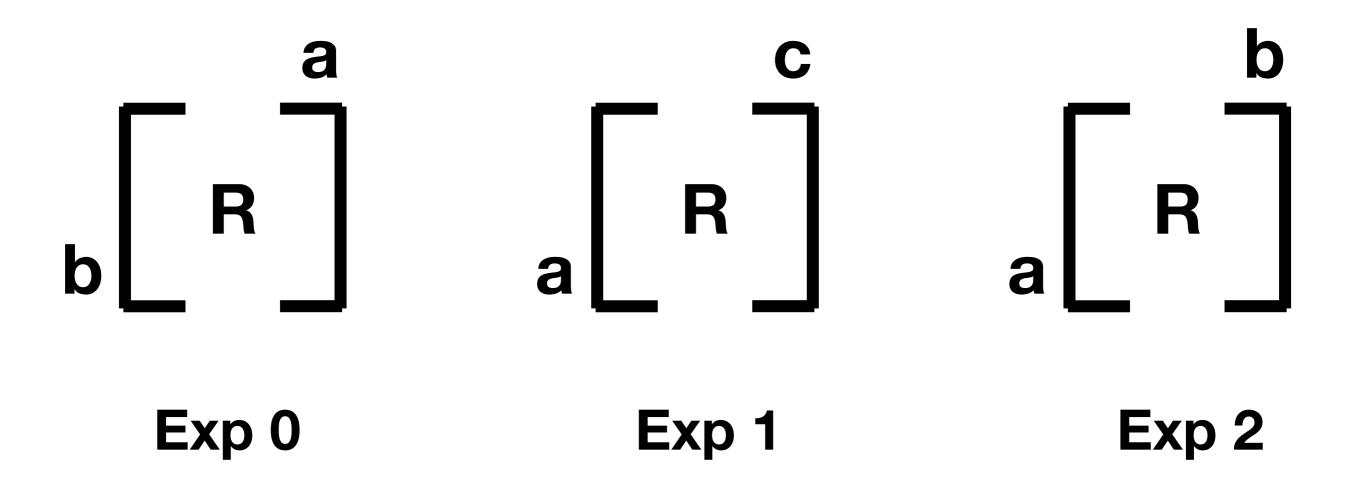
## Remove path from Query Graph



### Algebraic expression from path (chain)



### Expressions



### Arrangements

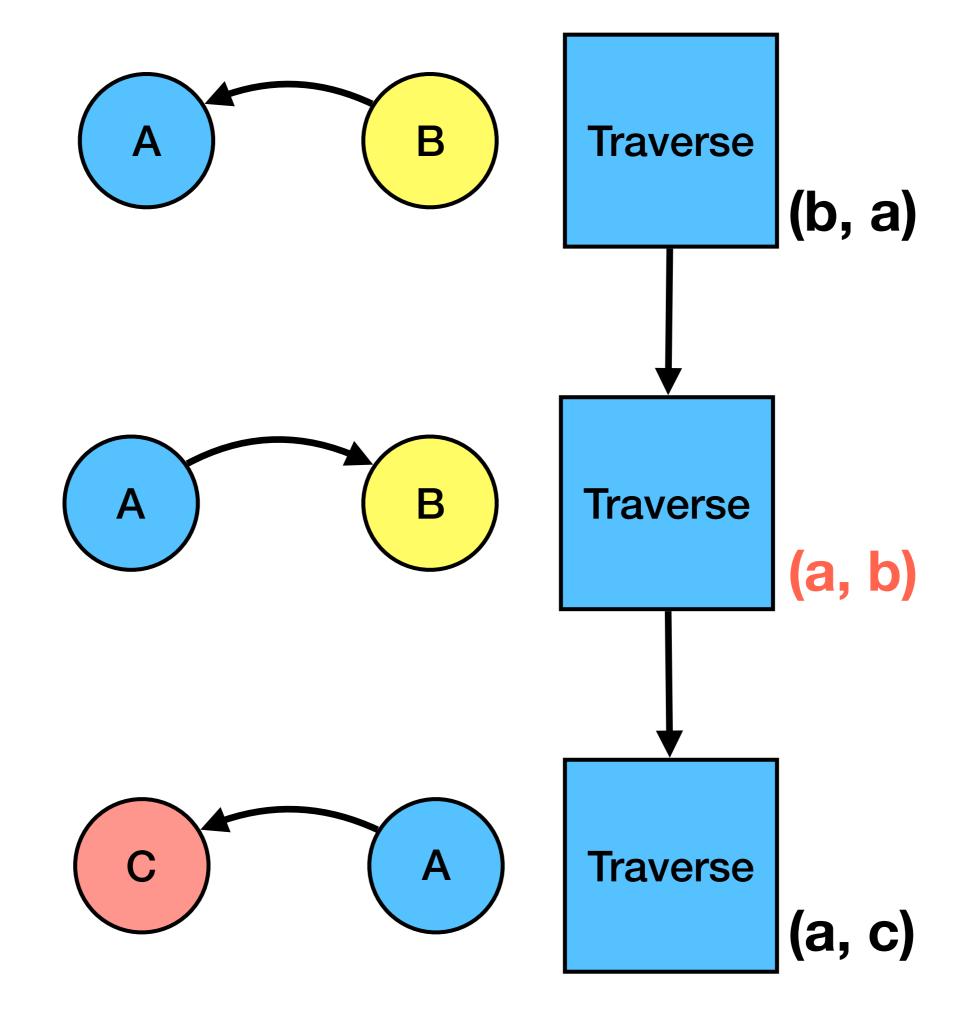
Permute ([Exp0, Exp1, Exp2]), N!

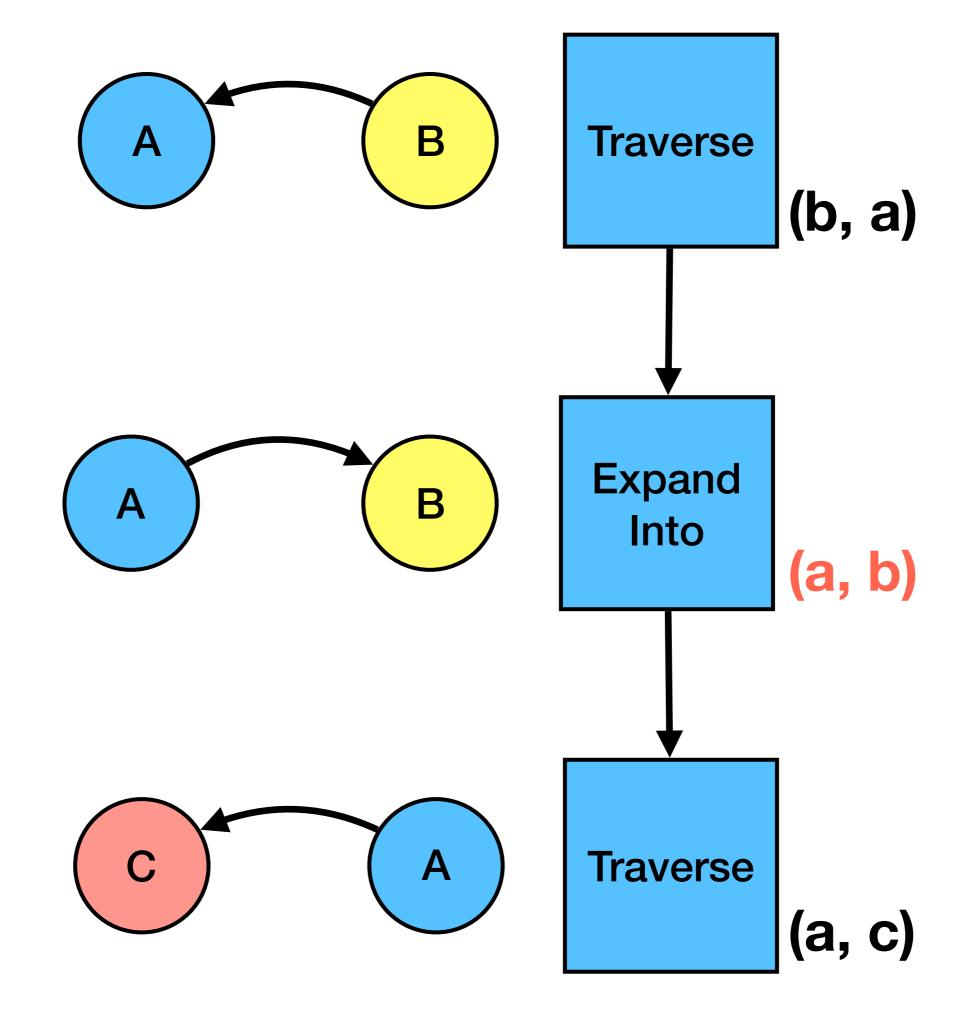
Pick best arrangement to be executed in order

Valid arrangement:

1. Expi either src or dest node is already resolved by previous expression.

2. Early filters





### Thank you

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